

Machine Programming Basics

CPSC 235 - Computer Organization

References

- Slides adapted from CMU

Outline

- History of Intel processors and architectures
- Assembly basics: registers, operands, move
- Arithmetic and logical operations
- C, assembly and machine code

Intel x86 Processors

- Dominate laptop/desktop/server market
- Evolutionary design
 - Backwards compatible up until 8086, introduced in 1978
 - Added more features as time goes on
- Complex instruction set computer (CISC)
 - Many different instructions with many different formats
 - Difficult to match performance of Reduced Instruction Set Computers (RISC)
 - But, Intel has done just that in terms of speed, less so for low power

Intel x86 Evolution: Milestones

Name	Date	Transistors	MHz	Notes
8086	1978	29K	5-10	16-bit
386	1985	275K	16-33	32-bit
Pentium 4E	2004	125M	2800-3800	64-bit
Core 2	2006	291M	1060-3333	multi-core
Core i7	2008	731M	1600-4400	four cores

x86 Clones: Advanced Micro Devices (AMD)

- Historically
 - AMD has followed just behind Intel
 - A little bit slower, a lot cheaper
- Then
 - Recruited top circuit designers from Digital Equipment Corp. and other downward trending companies
 - Built Opteron: tough competitor to Pentium 4
 - Developed x86-64, their own extension to 64 bits
- Recent years
 - Intel leads the world in semiconductor technology
 - AMD has fallen behind

Intel's 64 bit History

- 2001: Intel attempts radical shift from IA32 to IA64
 - Totally different architecture (Itanium)
 - Performance disappointing
- 2003: AMD steps in with evolutionary solution
 - x86-64 (now called "AMD64")
- 2004: Intel Announces EM64T extension to IA32
 - Extended Memory 64 bit Technology
 - Almost identical to x86-64
- All but low-end x86 processors support x86-64
 - but, lots of code still runs in 32 bit mode

Definitions

- Architecture: the parts of a processor design that one needs to understand for writing correct machine/assembly code
 - Machine code: the byte level programs that a processor executes
 - Assembly code: a text representation of machine code
- Microarchitecture: implementation of the architecture
- Example Instruction Set Architectures (ISA)
 - Intel: x86, IA32, Itanium, x86-64
 - ARM: Used in almost all mobile phones
 - RISC V: new open source ISA

Assembly/Machine Code View

- Programmer Visible State
 - PC: Program counter
 - Address of next instruction
 - Register file
 - Condition codes
 - store status information about most recent arithmetic or logical operation
- Memory
 - Byte addressable array
 - Code and user data
 - Stack to support procedures

Assembly Characteristics

- “Integer” data of 1, 2, 4, or 8 bytes
 - data values
 - addresses (untyped pointers)
- Floating point data of 4, 8, or 10 bytes
- SIMD vector data types of 8, 16, 32, or 64 bytes
- Code: byte sequences encoding series of instructions
- No aggregate types such as arrays or structures

x86-64 Integer Registers

8-byte register	bytes 0-3	bytes 0-1	byte 0
%rax	%eax	%ax	%al
%rcx	%ecx	%cx	%cl
%rdx	%edx	%dx	%dl
%rbx	%ebx	%bx	%bl
%rsi	%esi	%si	%sil
%rdi	%edi	%di	%dil
%rsp	%esp	%sp	%spl
%rbp	%ebp	%bp	%bpl

x86-64 Integer Registers (continued)

8-byte register	bytes 0-3	bytes 0-1	byte 0
%r8	%r8d	%r8w	%r8b
%r9	%r9d	%r9w	%r9b
%r10	%r10d	%r10w	%r10b
%r11	%r11d	%r11w	%r11b
%r12	%r12d	%r12w	%r12b
%r13	%r13d	%r13w	%r13b
%r14	%r14d	%r14w	%r14b
%r15	%r15d	%r15w	%r15b

x86-64 Integer Registers (continued)

- Some assembly instructions include a suffix that indicates what portion of the register is accessed:
 - **q**: “quadword” 8 bytes
 - **l**: “double word” lower 4 bytes
 - **w**: “word” lower 2 bytes
 - **b**: “byte” lowest byte

Assembly Characteristics: Operations

- Transfer data between memory and register
 - Load data from memory into register
 - Store register data into memory
- Perform arithmetic function on register or memory data
- Transfer control
 - Unconditional jumps to/from procedures
 - Conditional branches
 - Indirect branches

Moving Data

- Instruction:
 - `movq` *source* (Src), *destination* (Dest)
- Operand types
 - Immediate (Imm): constant integer data
 - Register (Reg): one of 16 integer registers
 - Memory (Mem): 8 consecutive bytes of memory at address given by register

movq Operand Combinations

Source	Destination	Example	C Analog
Imm	Reg	<code>movq \$0x4, %rax</code>	<code>temp = 0x04;</code>
Imm	Mem	<code>movq \$-147, (%rax)</code>	<code>*p = -147;</code>
Reg	Reg	<code>movq %rax, %rdx</code>	<code>temp2 = temp1;</code>
Reg	Mem	<code>movq %rax, (%rdx)</code>	<code>*p = temp;</code>
Mem	Reg	<code>movq (%rax), %rdx</code>	<code>temp = *p;</code>

Memory Addressing Modes

- Immediate

- `$val`
- `val`: constant integer value
- example: `movq $7, %rax`

- Normal

- `(R) Mem[Reg[R]]`
- `R`: register `R` specifies memory address
- `movq (%rcx), %rax`

Memory Addressing Modes (continued)

- Displacement

- $D(R) \text{ Mem}[\text{Reg}[R] + D]$
- R: register specifies start of memory region
- D: constant displacement D specifies offset
- example: `movq 8(%rdi), %rdx`

Memory Addressing Modes (continued)

- Indexed

- $D(Rb, Ri, S) \text{ Mem}[\text{Reg}[Rb] + S * \text{Reg}[Ri] + D]$
- D: constant displacement 1, 2, or 4 bytes
- Rb: base register
- Ri: index register: any except %esp
- S: scale: 1, 2, 4, or 8
- example: `movq 0x100(%rcx, %rax, 4), %rdx`

Addressing Modes Example

- Example C code

```
void swap (long *xp, long *yp) {  
    long t0 = *xp;  
    long t1 = *yp;  
    *xp = t1;  
    *yp = t0;  
}
```

Addressing Modes Example

- x86 assembly version

```
# %rdi = xp
```

```
# %rsi = yp
```

```
swap:
```

```
movq    (%rdi), %rax # t0 = *xp
```

```
movq    (%rsi), %rdx # t1 = *yp
```

```
movq    %rdx, (%rdi) # *xp = t1
```

```
movq    %rax, (%rsi) # *yp = t0
```

```
ret
```

Address Computation Examples

- `rdx` contains `0xf000`
- `rcx` contains `0x0100`

Expression	Address Computation	Address
<code>0x8(%rdx)</code>	<code>0xf000 + 0x8</code>	<code>0xf008</code>
<code>(%rdx, %rcx)</code>	<code>0xf000 + 0x100</code>	<code>0xf100</code>
<code>(%rdx, %rcx, 4)</code>	<code>0xf000 + 4*0x100</code>	<code>0xf400</code>
<code>0x80(,%rdx,2)</code>	<code>2*0xf000 + 0x80</code>	<code>0x1e080</code>

Address Computation Instruction

- `leaq Src, Dest`
 - Load effective address of source into destination
- Uses
 - Computing addresses without a memory reference
 - Computing arithmetic expressions of the form $x + k * y$
- Example

```
long m12(long x) {  
    return x*12;  
}
```

```
leaq (%rdi, %rdi, 2), %rax # t = x+2*x  
salq $2, %rax
```

Some Arithmetic Operations

- Binary operators

addq	Src, Dest	Dest = Dest + Src
subq	Src, Dest	Dest = Dest - Src
imulq	Src, Dest	Dest = Dest * Src
salq	Src, Dest	Dest = Dest << Src
sarq	Src, Dest	Dest = Dest >> Src (arithmetic)
shrq	Src, Dest	Dest = Dest >> Src (logical)
xorq	Src, Dest	Dest = Dest ^ Src
andq	Src, Dest	Dest = Dest & Src
orq	Src, Dest	Dest = Dest Src

- Be careful of the argument order

Some Arithmetic Operations

- Unary operators

incq	Dest	$\text{Dest} = \text{Dest} + 1$
decq	Dest	$\text{Dest} = \text{Dest} - 1$
negq	Dest	$\text{Dest} = - \text{Dest}$
notq	Dest	$\text{Dest} = \sim \text{Dest}$

Arithmetic Expression Example

- C code

```
long arith (long x, long y, long z) {  
    long t1 = x+y;  
    long t2 = z+t1;  
    long t3 = x+4;  
    long t4 = y * 48;  
    long t5 = t3 + t4;  
    long rval = t2 + t5;  
    return rval;  
}
```

Arithmetic Expression Example

- Assembly code

```
# %rdi = x  
# %rsi = y  
# %rdx = z
```

arith:

```
leaq    (%rsi,%rsi,2), %rax  
salq    $4, %rax  
leaq    4(%rdi,%rax), %rax  
addq    %rsi, %rdi  
addq    %rdx, %rdi  
addq    %rdi, %rax  
ret
```

Turning C into Object Code

- Code in files `p1.c` and `p2.c`
- Compile with command: `gcc -Og p1.c p2.c -o p`
 - use basic optimizations (`-Og`)
 - put resulting binary in file `p`
- The above `gcc` command runs the following programs:
 - source text → `cpp` → compiler → assembler → linker

Assembly

- Compiling C to assembly: `gcc -Og -S <file>`
 - produces an assembly file `<file>.s`
- Disassembling Code: `objdump -d <file>`
 - useful tool for examining object code
 - analyzes bit pattern of series of instructions
 - produces approximate rendition of assembly code

Summary

- History of Intel processors and architectures
- C, assembly, machine code
 - new forms of visible state: program counter, registers, . . .
 - Compiler must transform language constructs into low level instruction sequences
- Assembly basics: registers, operands, move
- Arithmetic