CSC552 – Advanced UNIX Programming

Threads

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Printer Server



UNIX Threads

- Exists within process and uses process resources
- Independent flow of control
- Share resources with other threads
- Dies with parent process

Process Structure

- Memory map (virtual addresses)
- File descriptor table
- Signal descriptor table
- User IDs
- Current working directory

Thread Maintains Own

Registers

- Stack Pointer
- Program Counter
- Scheduling properties
- Set of pending and block signals
- Thread specific data

Process Structure



Image from "Multithreaded Programming with Pthreads by Lewis & Berg

Context–Switch

Process

- Registers
- Virtual memory
- Some other process state

Threads

• Registers



Three Threads Running Concurrently on One CPU

Image from "Multithreaded Programming with Pthreads by Lewis & Berg

Parallelism



Three Threads Running in Parallel on Three CPUs

Image from "Multithreaded Programming with Pthreads by Lewis & Berg

Performance

Platform	fork()			pthread_create()		
	real	user	sys	real	user	sys
IBM332 MHz 604e 4 CPUs/node 512 MB Memory	92.4	2.7	105.3	8.7	4.9	3.9
AIA 4.3						
IBM 375 MHz POWER3 16 CPUs/node 16 GB Memory AIX 5.1	173.6	13.9	172.1	9.6	3.8	6.7
INTEL 2.2 GHz Xeon CPU/node 2 GB Memory RedHat Linux 7.3	17.4	3.9	13.5	5.9	0.8	5.3

Suitable Tasks

- Block for long waits
- Use many CPU cycles
- Must respond to asynchronous events
- Lesser or greater importance
- Performed in parallel

Common Models

- Manager/worker
- Pipeline
- Peer

POSIX Threads

- pthreads
- Return
 - 0
 - Error number
- Thread ID
 - pthread_t
- Basic Functions
 - pthread_self()
 - pthread_equal()

Thread Creation

- pthread_create()
- Main single, default thread
- After a thread has been created, how do you know when it will be scheduled to run by the OS?

Create Example

}

- pthread_t threads[NUM_THREADS]; int rc, t;

What is wrong with this code fragment?

How can it be corrected?

More Examples

- threads/hello.c
- threads/hello_struct.c

Freeing space

- Return value
 - Return function
 - Free space
- Parameter
 - Do not reuse
 - Separate variable

threads/copymultiple.c
Page 429, Ex 12.13

Thread Termination

- Thread returns from starting routine
- Calls pthread_exit()
- Cancelled by another thread calling pthread_cancel()
- pthread_detach()

Joinable Threads

- Detachable vs. Nondetachable
- Suspend execution of calling thread until joined thread terminates
 pthread_join()
- Single thread exit
 pthread_exit()

Cancel a Thread

- > pthread_cancel()
- > pthread_setcanceltype()
- > pthread_testcancel()

User-level Threads



Kernel-level Threads



Hybrid Threads



Mutex

Mutual exclusion

Two states

- Locked
- Unlocked

Synchronization Critical sections Shared resources

Typical Sequence

- Create and initialize
- Several threads attempt to lock the mutex
- Only one succeeds
- The owner thread performs some set of actions
- The owner unlocks the mutex
- Another thread acquires the mutex and repeats the process
- Finally the mutex is destroyed

Mutex Functions

- Type: pthread_mutex_t
- pthread_mutex_init()
- pthread_mutex_destroy()
- pthread_mutex_lock()
- pthread_mutex_trylock()
- pthread_mutex_unlock()

Cooperation

An uncooperative thread can enter a critical section without acquiring a mutex lock.

What are some ways to prevent this from happening?

Example

threads/counter.c

- What are possible side affects or what can go wrong if the *count* variable is not protected by a mutes lock?
- What really happens when a variable is incremented?

Protect Library Functions

- int randsafe(double *ranp) {
 static pthread_mutex_t lock =
 PTHREAD_MUTEX_INITIALIZER;
 int error;
 - if (error = pthread_mutex_lock(&lock))
 return error;
 *ranp = (rand() + 0.5)/(RAND_MAX + 1.0);
 return pthread_mutex_unlock(&lock);

Wait for Condition

- Busy wait
- Mutex
 - Lock mutex
 - Test condition
 - If true, unlock mutex and exit loop
 - If false, suspend thread and unlock mutex

Condition Variables

- Associated with specific condition
- Atomic waiting operation
- Type: pthread_cond_t
- Initialize
 - PTHREAD_CONDITION_INITIALIZER
 - o pthread_cond_init()
- Destroy
 - o pthread_cond_destroy()

Condition Wait

- pthread_cond_wait()
- > pthread_mutex_lock(&m); while (x != y) pthread_cond_wait(&v, &m); /* modify x or y if necessary */ pthread_mutex_unlock(&m);

Condition Signal

- pthread_cond_signal()
- pthread_cond_broadcast()
- pthread_mutex_lock(&m) x++; pthread_cond_signal(&v); pthread_mutex_unlock(&m);

pthread_cond_timedwait()

Guidelines

- Acquire the mutex before testing the predicate
- Retest the predicate after returning from a pthread_cond_wait - Why?
- Acquire the mutex before changing any of the variables appearing in the predicate.
- Hold the mutex only for a short period of
- Release the mutex
 - explicitly pthread_mutex_unlock()
 - implicitly pthread_cond_wait()

Examples

- threads/condvar1.c
- threads/tbarrier.c
- threads/syncConditionVar.c

Signal Handling and Threads

- All threads in process share process's signal handlers.
- Each thread has its own signal mask.

What does this mean?

Signals and Threads

- Signal Types
 - Asynchronous
 - Synchronous
 - Directed
- > pthread_sigmask()